



C340 Concurrency: Model-based Design

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Outline

- ***Role of Modelling in System Development***
- ***Refining Models into Designs***
 - ***FSP Actions and Operations***
 - ***FSP Processes and Threads***
 - ***FSP Processes and Monitors***
 - ***FSP Safety Properties and Monitor Invariants***
- ***Single-Lane Bridge Revisited***

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What are models?

- *Abstract representations of the real world*
- *Neglect details considered unnecessary*
- *Example: mathematical model of bridge*
- *Supports analysis to check properties of interest*

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Why do we model?

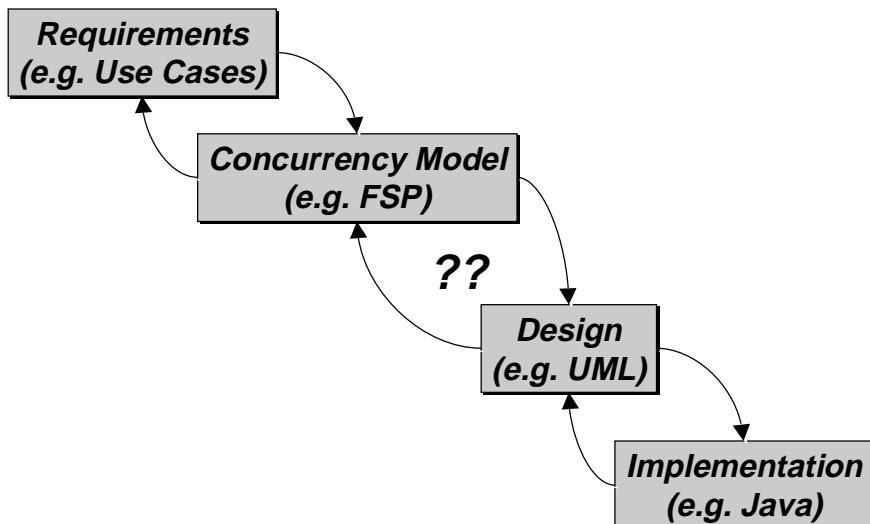
- *Guide system design*
- *Improve system design in a cost-effective way*
- *Obtain formal basis for verifying correctness of the system*

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Engineering Concurrent Programs



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Model-based Design

- ***What is the relationship between concurrency model (e.g. in FSP) and design (e.g. in UML)?***
- ***How can we derive the UML design for a concurrent program from an FSP model?***
- ***How do we transfer properties that have been proven to hold in FSP into an UML design?***

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Single Lane Bridge FSP



```
CAR = (request->enter->exit->CAR).
NOPASS1=C[1], C[i:ID]=((i).enter->C[i%N+1]).
NOPASS2=C[1], C[i:ID]=((i).exit->C[i%N+1]).
|| CONVOY=((ID):CAR|| NOPASS1|| NOPASS2).
|| CARS=(red:CONVOY|| blue:CONVOY).
SAFEBRIDGE=SAFEBRIDGE[0][0],
SAFEBRIDGE[nr:T][nb:T]=
  (when (nb==0) red[ID].enter->BRIDGE[nr+1][nb]
  |when (nr==0) blue[ID].enter->BRIDGE[nr][nb+1]
  |when (nr>0) red[ID].exit -> BRIDGE[nr-1][nb]
  |when (nb>0) blue[ID].exit-> BRIDGE[nr][nb-1]).
FAIRBRIDGE=FAIRBRIDGE[0][0][0][0][True],
FAIRBRIDGE[nr:T][nb:T][wr:T][wb:T][bt:B]=
  (when(wr<N) red[ID].request->BRIDGE[nr][nb][wr+1][wb][bt]
  |when(wb<N) blue[ID].request->BRIDGE[nr][nb][wr][wb+1][bt]
  |when(nb==0&&(wb==0 || !bt))red[ID].enter->BRIDGE[nr+1][nb][wr-1][wb][bt]
  |when(nr==0&&(wr==0 || bt))blue[ID].enter->BRIDGE[nr][nb+1][wr][wb-1][bt]
  |when(nr>0) red[ID].exit -> BRIDGE[nr-1][nb][wr][wb][True]
  |when(nb>0) blue[ID].exit-> BRIDGE[nr][nb-1][wr][wb][False]).
|| SINGLELANEBRIDGE = (CARS || FAIRBRIDGE)>>{red[ID].exit,blue[ID].exit}.
property ONEWAY=(red[ID].enter -> RED[1] | blue[ID].enter ->BLUE[1],
               RED[i:ID]=(red[ID].enter->RED[i+1]
                           |when (i==1)red[ID].exit->ONEWAY)
                           |when (i>1) red[ID].exit->RED[i-1]),
               BLUE[i:ID]=(blue[ID].enter->BLUE[i+1]
                           |when (i==1)blue[ID].exit->ONEWAY)
                           |when (i>1) red[ID].exit->BLUE[i-1]).
|| TESTSAFETY = (CARS || SAFEBRIDGE).
```

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Deriving Thread Classes

- *Those FSP processes that represent active entities*
- *Candidates: Processes that are instantiated multiple times*
- *Example in Single Lane Bridge: Cars*

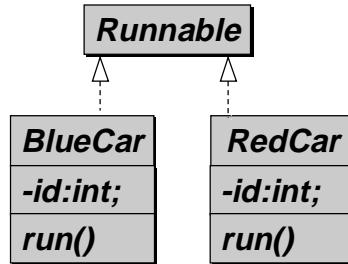
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Deriving Thread Classes

```
CAR = (request->enter->exit->CAR).  
NOPASS1=C[1], C[i:ID]=([i].enter->C[i%N+1]).  
NOPASS2=C[1], C[i:ID]=([i].exit->C[i%N+1]).  
|| CONVOY=([ ID]:CAR | NOPASS1 | NOPASS2).  
|| CARS=(red:CONVOY | blue:CONVOY).
```



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Deriving Monitors

- *Those FSP processes that are not active do not have to be separate threads*
- *If they represent shared resources are likely to become monitors*
- *Example in FSP: Bridge*

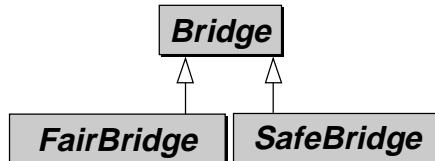
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Deriving Monitors

FAIRBRIDGE[nr:T][nb:T][wr:T][wb:T][bt:B]= ...
SAFEBRIDGE[nr:T][nb:T]= ...



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Deriving Monitor Variables

- *FSP Processes that represent monitors often maintain a local state in terms of index variables*
- *These index variables also have to be used monitors*
- *They have to become instance variables of the monitor*
- *Example: Instance Variables of FairBridge and SafeBridge*

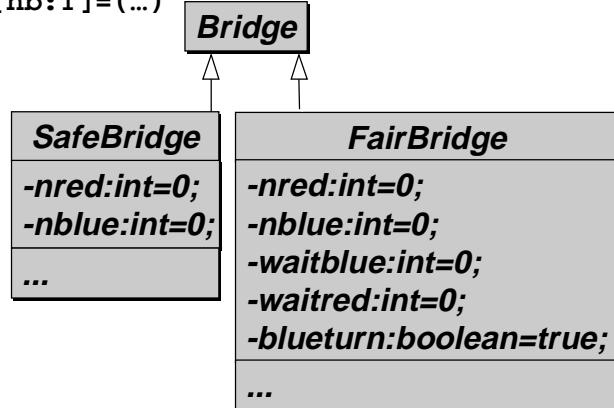
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Deriving Monitor Variables

```
FAIRBRIDGE=FAIRBRIDGE[0][0][0][0][True],  
FAIRBRIDGE[nr:T][nb:T][wr:T][wb:T][bt:B]=(...).  
SAFEBRIDGE=SAFEBRIDGE[0][0],  
SAFEBRIDGE[nr:T][nb:T]=(...)
```



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Deriving Monitor Operations

- *FSP processes that became monitors engage in actions*
- *These actions are candidates to become operations in monitor classes*
- *For implementation purposes it may be appropriate to subsume operations or split these operations*

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Deriving Monitor Operations

```
FAIRBRIDGE[nr:T][nb:T][wr:T][wb:T][bt:B]=  
  (when(wr<N) red[ID].request->BRIDGE[nr][nb][wr+1][wb][bt]  
  |when(wb<N) blue[ID].request->BRIDGE[nr][nb][wr][wb+1][bt]  
  |when(nb==0&&(wb==0||!bt))red[ID].enter->BRIDGE[nr+1][nb][wr-1][wb][bt]  
  |when(nr==0&&(wr==0||bt)) blue[ID].enter->BRIDGE[nr][nb+1][wr][wb-1][bt]  
  |when(nr>0) red[ID].exit -> BRIDGE[nr-1][nb][wr][wb][True]  
  |when(nb>0) blue[ID].exit-> BRIDGE[nr][nb-1][wr][wb][False]).
```

```
SAFEBRIDGE[nr:T][nb:T]=
```

```
  (when (nb==0) red[ID].enter->BRIDGE[nr+1][nb]  
  |when (nr==0) blue[ID].enter->BRIDGE[nr][nb+1]  
  |when (nr>0) red[ID].exit -> BRIDGE[nr-1][nb]  
  |when (nb>0) blue[ID].exit-> BRIDGE[nr][nb-1]).
```

Bridge

| FairBridge | SafeBridge |
|--|--|
| ... | ... |
| <i>redEnter()</i> <i>redExit()</i> <i>blueEnter()</i> <i>blueExit()</i> | <i>redEnter()</i> <i>redExit()</i> <i>blueEnter()</i> <i>blueExit()</i> |

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Deriving Monitor Conditions

- Guards of those FSP processes that have become monitors determine conditions for synchronization
- Can be directly derived by negation
- Negation due to opposite perspectives:
 - Guard: When can action happen?
 - Monitor: How long do we have to wait before action can happen?

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Deriving Monitor Conditions

■ FSP:

```
FAIRBRIDGE[nr:T][nb:T][wr:T][wb:T][bt:B]=  
...  
|when(nb==0&&(wb==0||!bt))red[ID].enter->...
```

■ Java:

```
class FairBridge {  
    synchronized void redEnter() throws InterruptedException{  
        ++waitred;  
        while (nblue>0 || (waitblue>0 && blueturn)) wait();  
        --waitred;  
        ++nred;  
    }  
    ...  
}
```

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Deriving Monitor Invariants

- **Safety properties that affect FSP processes that have become monitors have to hold all the time**
- **Safety properties become monitor invariants**

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Deriving Monitor Invariants

■ **FSP:**

```
property ONEWAY=(red[ID].enter -> RED[1]
                | blue[ID].enter ->BLUE[1],
                RED[i:ID]=(red[ID].enter->RED[i+1]
                            |when (i==1)red[ID].exit->ONEWAY)
                            |when (i>1) red[ID].exit->RED[i-1]),
                BLUE[i:ID]=(blue[ID].enter->BLUE[i+1]
                            |when (i==1)blue[ID].exit->ONEWAY)
                            |when (i>1) red[ID].exit->BLUE[i-1]).
```

■ **Java:**

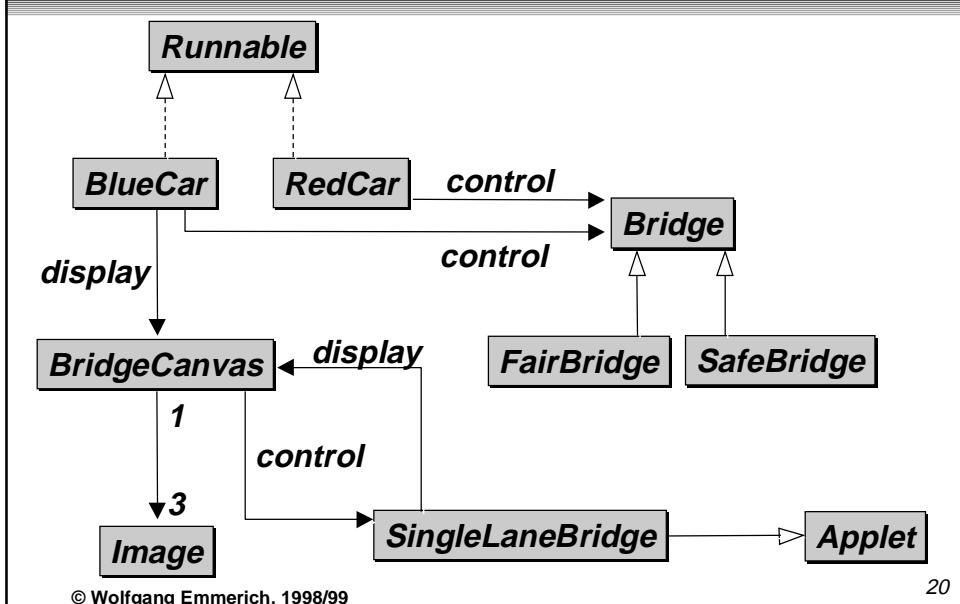
```
class SafeBridge extends Bridge {
    synchronized void redEnter() throws InterruptedException
    { while (nblue>0) wait(); ++nred; }
    synchronized void redExit()
    { --nred; if (nred==0) notifyAll();}
    synchronized void blueEnter()throws InterruptedException
    { while (nred>0) wait();++nblue; }
    synchronized void blueExit()
    { --nblue; if (nblue==0) notifyAll(); }
}
```

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UML Overview of Single Lane Bridge



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Summary

- ***Role of Modelling in System Development***
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