



C340 Concurrency: Message Passing

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Outline

- ***Motivation***
- ***Synchronous Message Passing***
- ***Modelling Synchronous Message Passing
in FSP***
- ***Selective Receive***
- ***Asynchronous Message Passing***
- ***Modelling Asynchronous Message
Passing in FSP***
- ***Rendezvous in Java and FSP***



Absence of Shared Memory

- ***In previous lectures interaction between threads via shared memory***
- ***In java references to shared objects***
- ***Usually encapsulated in Monitors***
- ***In a distributed setting shared memory does not exist***
- ***Communication is achieved via passing messages between concurrent/parallel threads***

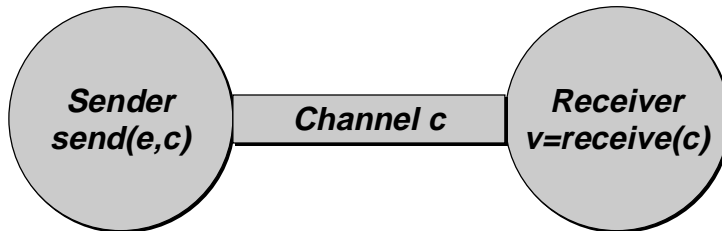


Forms of Message Passing

- ***Principal Operations***
 - *send*
 - *receive*
- ***Synchronization***
 - *Synchronous*
 - *Asynchronous*
 - *Rendevous*
- ***Multiplicity***
 - *one-one*
 - *many-one*
 - *many-many*
- ***Anonymity***
 - *anonymous message passing*
 - *non-anonymous message passing*
- ***Receipt of Messages***
 - *Full*
 - *Selective*



Synchronous Message Passing



- *send(e,c): Send e to channel c. Sending process is blocked until channel received e*
- *v=receive(c): Receive a value into a local variable v from channel C. The calling process is blocked until a message is sent into channel*
- *No buffering for synchronous messages*

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Sync. Message Passing in Java

- *Encapsulate message passing abstractions in monitor Channel:*

```
class Channel extends Selectable {
    public synchronized void send (object v)
        throws InterruptedException {...}
    public synchronized Object receive() {...}
}
```
- *What are the synchronization conditions?*

Demo

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Modelling Sync. Message Passing

```

range M=0..9
SENDER = SENDER[0],
SENDER[e:M]=(chan.send[e]->SENDER[(e+1)%10]).
RECEIVER = (chan.receive[v:M]->RECEIVER).
|| SYNCMSG = (SENDER || RECEIVER)
               /{chan/chan.{send,receive}}.

```

■ To avoid re-labelling:

```

range M=0..9
SENDER = SENDER[0],
SENDER[e:M]=(chan[e]->SENDER[(e+1)%10]).
RECEIVER = (chan[v:M]->RECEIVER).
|| SYNCMSG = (SENDER || RECEIVER).

```

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Selective Receives

■ Receiving from more than one channel

■ Stuck if we choose the wrong channel

■ Selective receives (e.g. Occam or Ada):

```

select when G1 and v1=receive(chan1) => S1;
      or when G2 and v2=receive(chan2) => S2;
      ...
      or when G3 and vn=receive(chann) => S3;
end;

```

■ Note similarity to FSP guarded actions

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Modelling Selective Receives

■ Example: Car Park Control

```
CONTROL(N=4) = SPACES[N],  
SPACES[i:0..N]=(when(i>0)arrive->SPACES[i-1]  
                |when(i<N)depart->SPACES[i+1]  
                ).
```

```
ARRIVALS=(arrive->ARRIVALS).
```

```
DEPARTURES=(depart->DEPARTURES).
```

```
||CARPARK=(ARRIVALS || DEPARTURES || CONTROL(4)).
```

■ How to implement CONTROL using Message Passing?

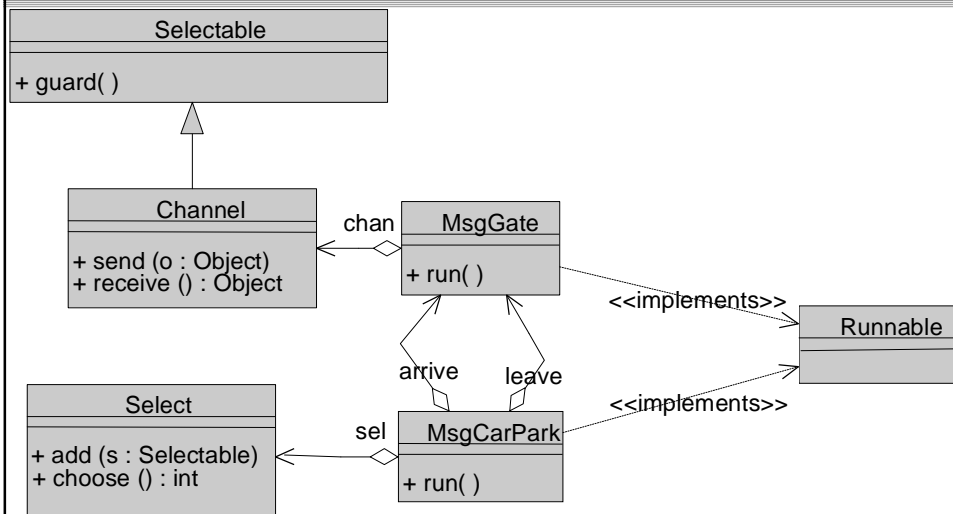
Demo

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Classes for Selective Receive



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Implementing Selective Receives

```
class MsgGate implements Runnable {
    private Channel chan;
    private Object signal = new Object();
    public MsgGate (Channel c) {chan=c;}
    public void run() {
        try {
            while(true) {
                ThreadPanel.rotate(12);
                chan.send(signal);
                ThreadPanel.rotate(348);
            }
        } catch (InterruptedException e){}
    }
}
```

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Implementing Selective Receives

```
class MsgCarPark implements Runnable {
    private Channel arrive,leave;
    private int spaces,N;
    public MsgCarPark(Channel a, Channel l, int capacity) {
        leave=l; arrive=a; N=spaces=capacity;
    }
    public void run() {
        try {
            Select sel = new Select();
            sel.add(leave); sel.add(arrive);
            while(true) {
                ThreadPanel.rotate(12);
                arrive.guard(spaces>0);
                leave.guard(spaces<N);
                switch (sel.choose()) {
                    case 1:leave.receive();display(++spaces); break;
                    case 2:arrive.receive();display(--spaces); break;
                }
            }
        } catch (InterruptedException e){}
    }
}
```

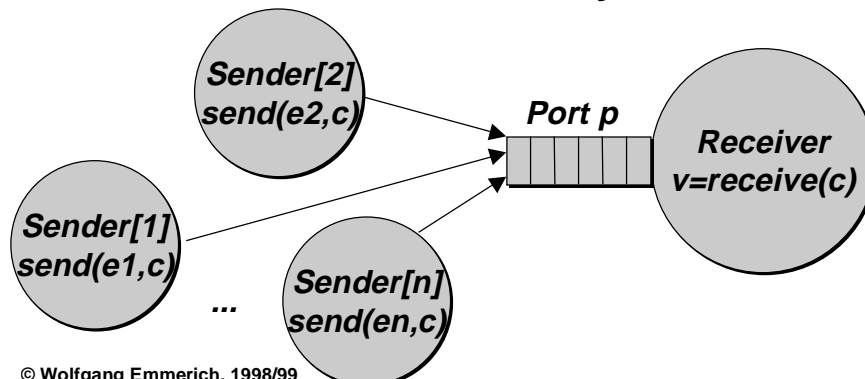
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Asynchronous Message Passing

- *Send does not block*
- *Messages are queued at the receiver*
- *We refer to these queues as ports*
- *Communication can be many-to-one*



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Async. Message Passing in Java

- *Two operations*
 - *send(e,p): send value e to port p. Calling process not blocked*
 - *v=receive(p): receive value into var v from port p. Calling process is blocked if no value queued to port.*

- *Implementation of Ports in Java:*

```
class Port extends Selectable{
    Vector queue;
    public synchronized void send(Object v) {...}
    public synchronized Object receive()
        throws InterruptedException {...}
}
```

Demo

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Modelling Async. Message Passing

```

range M = 0..4
set S = {[M],[M][M]}
PORT = (send[x:M]->PORT[x]),
PORT[h:M] = (send[x:M] ->PORT[x][h]
              |receive[h]->PORT),
PORT[t:S][h:M] = (send[x:M] ->PORT[x][t][h]
                  |receive[h]->PORT[t]).
| | APORT=PORT/{send/send[M],receive/receive[M]}.
ASENDER=ASENDER[0],
ASENDER[e:M]=(port.send[e]->ASENDER[(e+1)%4]).
ARECEIVER=(port.receive[v:M]->ARECEIVER).
| | ASYNCM=(s[1..2]:ASENDER | port:PORT | ARECEIVER)
  /{s[1..2].port.send/port.send}.

```

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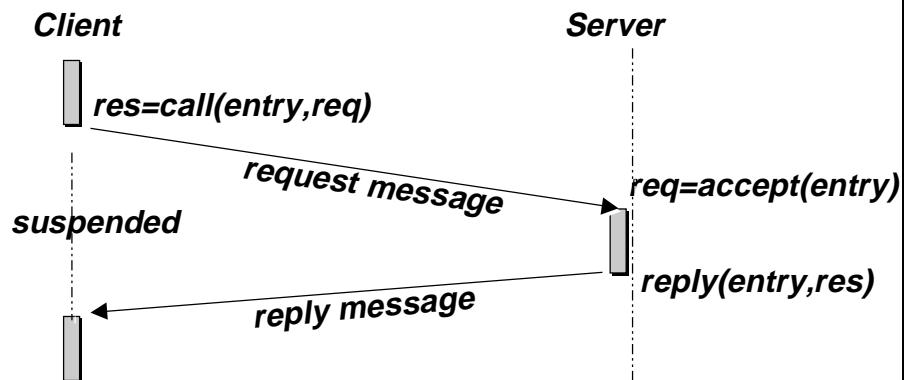
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Rendezvous Message Passing

- Request-Reply Protocol to support client-server interaction

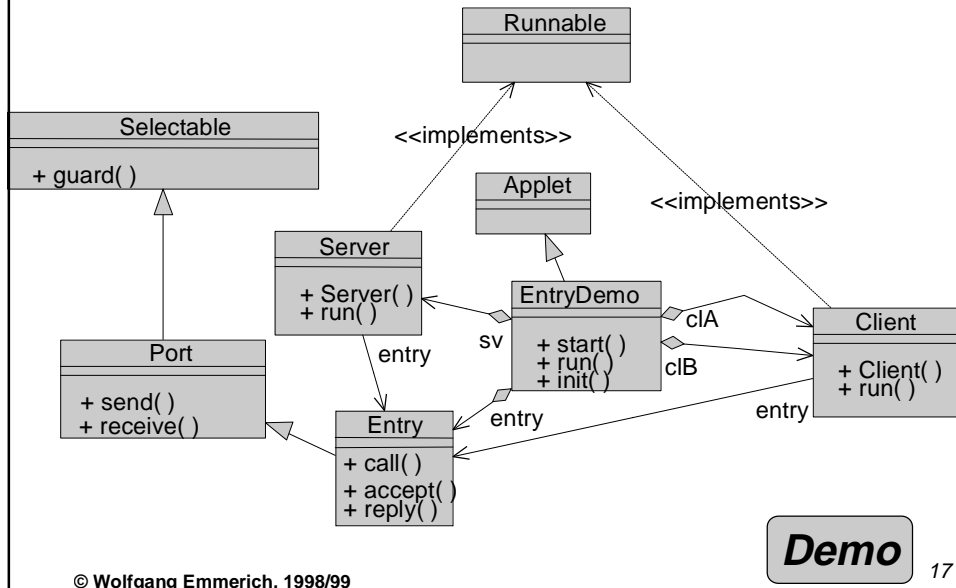


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Rendevous in Java



Summary

- **Synchronous Message Passing**
- **Modelling Synchronous Message Passing in FSP**
- **Selective Receive**
- **Asynchronous Message Passing**
- **Modelling Asynchronous Message Passing in FSP**
- **Rendevous**