

The TAOS Authentication System: Reasoning Formally About Security

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Motivation: Building Correct Authentication Systems

- We've studied **cryptographic primitives**
- We've studied **certificates**, and how they're used in SSL
 - Trusted third party, CA, attests to binding between public key and principal's name
 - One party can authenticate other using certificate
- Certificates are **more general tool**, but can be **hard to reason about**
- How can we **reason formally** about whether collection of certificates truly authenticates some principal to complete some operation on some object?

Motivation:

Flexible Authentication Systems

- Suppose want to authenticate user on client workstation to file server
 - User is principal
 - User authorized on file server to perform certain operations on certain file objects
- Simple model:
 - Use public-key cryptography
 - Install user's public key on file server
 - User holds private key on client workstation while logged in
 - User signs each RPC sent to file server using his private key

Motivation: Drawbacks of Simple Authentication Model

- **Very slow** (TAOS took 250 ms per RSA sig)
- **Rigid:**
 - What if I ssh into second machine?
 - 2nd box must sign RPCs to file server, too
 - Does it send messages back to 1st box for signing? **How would user know they're authentic?**
 - What if user goes home, leaves simulation running for hours?

Motivation: SSL Rigid, Too

- Does SSL work here?
- Assume both sides (client and server) authenticate by presenting certificates
- Fast: symmetric-key ciphers for session data
- But workstation must hold private key for every connection
- What if I ssh into second machine?
 - Want it to be able to use file server, too
 - Would have to give second machine my private key!

Outline of TAOS Authentication (1)

- Give each machine an identity:
public/private key pair
- User bkarp logs into machine X, **signs certificate**:
 - “bkarp **says** X **speaks for** bkarp.”
 - Reflects reality; X executes bkarp’s programs
 - In paper, **speaks for** written as \Rightarrow
 - Y **says** X just means “Y signs statement X with K_Y ” (note paper refers to **public key** when signing!)

Outline of TAOS Authentication (2)

- Now machine X can:
 - Open SSL-like secure channel from self to server; file server knows it's talking to X
 - Present “bkarp **says X speaks for** bkarp” to file server; file server knows X can speak for user
 - Send RPCs generated by bkarp's programs to file servers
 - All without machine X holding bkarp's private key!

Authorizing 2nd Machine with TAOS

- Consider ssh by bkarp to 2nd machine
- Want Y to talk to file server for bkarp
- ssh on X signs "X says Y can speak for bkarp"
- Gives this certificate to Y when bkarp logs into Y
- Now Y presents proof outline to file server:
 - I'm Y
 - X says Y can speak for bkarp
 - bkarp says X can speak for bkarp
- File server can check signatures and verify that RPCs authorized!

Why Can't SSL Authorize 2nd Machine?

- SSL for exactly two principals, tied to channels
- If X says something to Y, Y can't prove anything to Z
- In fact, Y can't verify anything after X closes its connection to Y
- SSL too rigid to support distributed systems with > 2 parties

TAOS's Central Strengths

- Certificates are true independent of channels
- ...so can be stored, passed to other parties
- ...and used to prove transitive trust relationships

Axioms in the TAOS Logic (2.1 in paper)

- speaks for:
 - if (A **speaks for** B) and (A **says** S)
then (B **says** S)
- handoff axiom:
 - if A **says** (B **speaks for** A)
then (B **speaks for** A)
- delegation axiom:
 - if A **says** (B | A) **speaks for** (B for A))
then (B | A) **speaks for** (B for A))

Applying Handoff and Delegation

- Handoff: given
A **says** (B **speaks for** A) and B **says** S
then A **says** S
- Delegation: given
A **says** (B | A) **speaks for** (B **for** A) and
B **says** A **says** S
then (B **for** A) **says** S

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Delegation more specific than handoff; records both principals, the trustor and trustee

Better for auditing...

Using Logic to Reason About Authentication

- Consider example in Section 2.2 of TAOS paper:
 - User Bob logs into workstation WS
 - Logic used to authenticate requests from Bob's login session to a remote file server FS
- What principals are involved?
 - Workstation firmware, OS, Bob, Channel
- Keep track of who knows:
 - Private keys
 - Signed certificates
 - Channel keys

State Before Bob Logs In

- Workstation firmware knows K_{vax4}
- User knows K_{bob} 's private "half"
- File server has K_{bob} 's public "half" in an ACL

Workstation Boot Time: Generating K_{ws}

- At boot, workstation firmware generates fresh public/private key, K_{ws}
- **Why not just use K_{vax4} directly?**
 - Don't want it to be stolen
 - Don't want statements to survive reboot (i.e., certificates generated for login sessions)
- Firmware signs:
“ K_{vax4} **says** (K_{ws} **speaks for** K_{vax4})”
- K_{vax4} never used again (until reboot)
- Why bother preserving K_{vax4} 's identity?
 - Why not just use K_{ws} as workstation's true identity?
 - **Want workstation's identity to survive reboots**

Boot Time: Generating K_{ws} (2)

- Why bother with roles (" K_{vax4} **as** OS")?
 - User might not trust some versions of OS, or some OS
 - Want to allow OS type/version to be visible in ACLs
 - Assuming a role amounts to reducing access rights
- Now vax4's authentication agent knows:
 - K_{ws} (but forgets K_{vax4})
 - $(K_{vax4}$ **as** OS) **says** (K_{ws} **speaks for** (K_{vax4} **as** OS))
- Why does vax4 need an identity at all?
 - So Bob can delegate to it!

Login: Delegation of Authority to Workstation by User

- Want ws to be able to act for Bob
- Bob signs with his private key, K_{bob} :
 K_{bob} **says** $((K_{\text{ws}} \mid K_{\text{bob}})$ **speaks for** $(K_{\text{ws}}$ **for** $K_{\text{bob}}))$
- Private half of K_{bob} not used again until next login!
- Why not “ K_{bob} **says** $(K_{\text{ws}}$ **speaks for** $K_{\text{bob}})$ ”?
 - If K_{ws} signs something, on whose behalf was it?
 - So statements by K_{ws} **ambiguous**, and perhaps **usable out of context**

Delegation at Login (2)

- What does $(A \mid B)$ mean?
 - That A is doing the signing
 - That A is claiming (no proof yet) that A is speaking for B
 - Really means that A says in its signed statement that it's speaking for B
- What does $(A \text{ for } B)$ mean?
 - Logical conclusion that A allowed to speak for B
 - i.e., $(A \mid B)$ plus delegation, like one on previous slide (see delegation axiom on p. 4 of paper)
 - By default, interpreted as B for purposes of ACLs
 - But for those who care, preserves who actually signed (A)

Delegation at Login (3)

- After delegation by Bob, vax4's authentication agent knows:

K_{ws} private half

$(K_{vax4} \text{ as OS}) \text{ says } (K_{ws} \text{ speaks for } (K_{vax4} \text{ as OS}))$

$K_{bob} \text{ says } ((K_{ws} \mid K_{bob}) \text{ speaks for } (K_{ws} \text{ for } K_{bob}))$

TAOS Channels

- TAOS uses symmetric-key ciphers to encrypt channels between hosts
- Channels named by their symmetric key
 - Name has global meaning
- C_{bob} **doesn't imply anything about Bob**
 - Only a mnemonic used by authors to indicate intent that C_{bob} carries messages from Bob
 - System must establish proof that this is case
- File server knows:
 - C_{bob} **says** RQ (where RQ a file server request)
 - i.e., "received request from someone who knows key C_{bob} "
- But **who** knows key C_{bob} ?
 - K_{ws} ?
 - K_{ws} on behalf of Bob?
 - K_{ws} on behalf of someone else?

Proving Authenticity: Channel Certificates

- ws signs **channel certificate** when channel between ws and file server first created:
 $(K_{ws} \mid K_{bob})$ **says** $(C_{bob}$ **speaks for** $(K_{ws}$ **for** $K_{bob}))$
- Goal: **link RQ encrypted with C_{bob} to Bob**
- Why not just have K_{bob} sign:
 - “ C_{bob} speaks for K_{bob} ”
 - This is what SSL client-side certificates do.
 - But in TAOS, **authentication agent doesn't hold K_{bob} 's private half**—and that's a good thing...

Channel Certificates (2):

- Why not have K_{ws} sign:
 - “ C_{bob} **speaks for** K_{ws} ”
 - Along with pre-signed “ K_{ws} **speaks for** K_{bob} ”
 - C_{bob} doesn't speak for K_{ws} in general! Only K_{bob} .
- Channel certificate is in fact **nicely restricted**:
 - States what we mean, and no more
 - vax4 **says** C_{bob} **speaks for** (vax4 **speaking for** Bob)
- But vax4 could sign this statement without Bob's agreement!
- So file server needs further evidence:
 - Is vax4 allowed to speak for Bob?

Using Logic to Prove Authenticity

- Suppose ws sends **all certificates** to file server:

$(K_{\text{vax4}} \text{ as OS})$ **says** $(K_{\text{ws}} \text{ speaks for } (K_{\text{vax4}} \text{ as OS}))$

K_{bob} **says** $((K_{\text{ws}} \mid K_{\text{bob}}) \text{ speaks for } (K_{\text{ws}} \text{ for } K_{\text{bob}}))$

$(K_{\text{ws}} \mid K_{\text{bob}})$ **says** $(C_{\text{bob}} \text{ speaks for } (K_{\text{ws}} \text{ for } K_{\text{bob}}))$

- Now file server can reason about meaning of C_{bob} **says** RQ

Using Logic to Prove Authenticity (2)

- File server can take
 K_{bob} **says** $((K_{\text{ws}} \mid K_{\text{bob}})$ **speaks for** $(K_{\text{ws}}$ **for** $K_{\text{bob}}))$
- and deduce, using **delegation axiom**:
 $(K_{\text{ws}} \mid K_{\text{bob}})$ **speaks for** $(K_{\text{ws}}$ **for** $K_{\text{bob}})$
- Informally, delegation axiom just means:
 - If Bob signs certificate allowing K_{ws} to speak for Bob, then K_{ws} is allowed to speak for Bob
- Really, delegation certificate means:
 - If K_{ws} says it's speaking for Bob, believe it.
 - This is **different than** " K_{ws} **speaks for** K_{bob} "!

Using Logic to Prove Authenticity (3)

- Now, combine:
 $(K_{ws} \mid K_{bob})$ **speaks for** $(K_{ws}$ **for** $K_{bob})$
 $(K_{ws} \mid K_{bob})$ **says** $(C_{bob}$ **speaks for** $(K_{ws}$ **for** $K_{bob}))$
- And thus derive:
 $(K_{ws}$ **for** $K_{bob})$ **says** $(C_{bob}$ **speaks for** $(K_{ws}$ **for** $K_{bob}))$
- In other words:
 - K_{ws} really does speak for K_{bob} ; it's not just claiming to do so
- So we can conclude that C_{bob} speaks for K_{ws} speaking for K_{bob}
- And thus:
 $(K_{ws}$ **for** $K_{bob})$ says RQ

TAOS: Summary

- Certificates allow flexible authentication
 - Can survive longer than a channel
 - Allow delegation of authority
 - Can be combined using formal logic
- Central ideas:
 - **says** and **speaks for**
 - handoff, delegation axioms
 - useful tools for reasoning formally about authentication in any distributed system!