

# **Introduction to Security and User Authentication**

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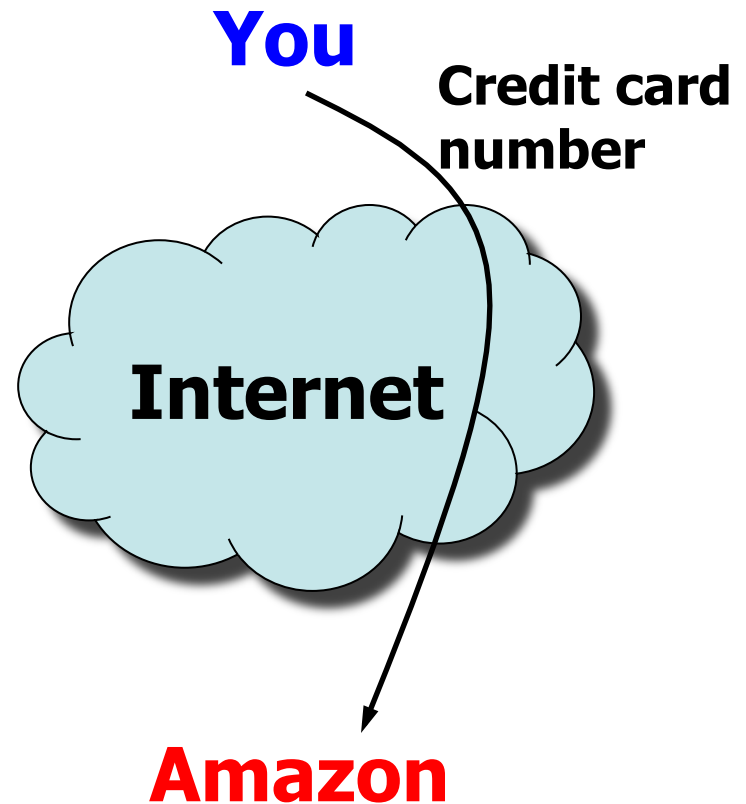
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# Topics We'll Cover

- User login authentication (local and remote)
- Cryptographic primitives, how to use them, and how **not** to use them
- Kerberos distributed authentication system
- Secure Sockets Layer (SSL)/Transport Layer Security (TLS) authentication and encryption system
- TAOS: logic for reasoning formally about authentication
- Software vulnerabilities and exploits
- Exploit Defenses
- Software Fault Isolation (SFI): containing untrusted code
- OKWS: a least-privilege isolated web server for UNIX

# A Simple Example

- Suppose you place an order with Amazon
- Goals:
  - You get the item you ordered
  - Amazon gets payment in the amount you agreed to pay on the payment page



# A Simple Example

- Suppose you place an order with Amazon

**How might this go wrong?  
Let us count the ways...**

- you ordered
- Amazon gets payment in the amount you agreed to pay on the payment page

**You**

**Credit card  
number**

**Amazon**

# Worries for Amazon Order

- What if an **eavesdropper taps Internet link?**
  - Network cables usually not physically secure
- What if someone has **broken into Internet router?** (They're just computers...)
- How do you know you're communicating with **Amazon?**
- How does Amazon know you are **authorized** to use the credit card number you provide?
- What if a **dishonest Amazon employee** learns my credit card number?
- What if Amazon sends me **wrong book, in error**

## Worries for Amazon Order (2)

- What if someone has broken into my desktop PC? Or my file server?
- Where did my web browser come from? How about my OS?
- What if my display or keyboard radiates a signal that can be detected at some distance?

# Worries for Amazon Order (2)

- What if someone has broken into my

**Fundamental security question:**

**“Whom or what am I trusting?”**

**Weakest item on list of answers determines system security!**

- What if my display or keyboard radiates a signal that can be detected at some distance?

# Whom or What Am I Trusting?

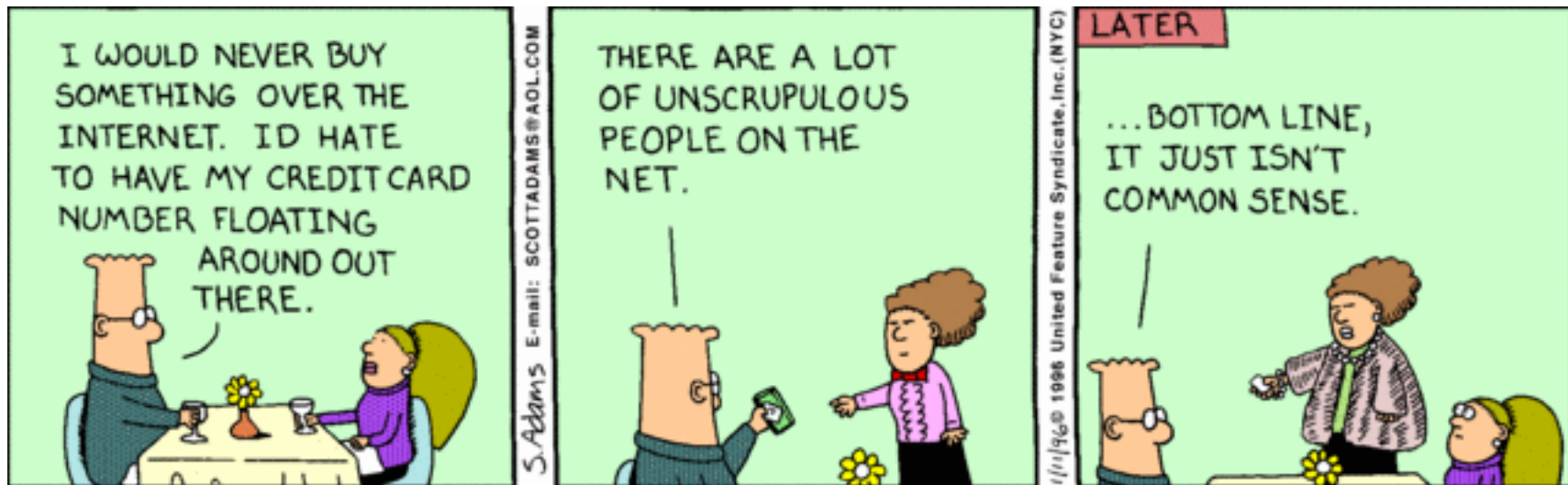
“They showed me a telephone, and said they were worried about ‘the microphone.’ When I look at a telephone, I see one high-fidelity microphone and one ‘low-fidelity microphone.’”

“Most people call this a telephone cord. I call it an antenna.”

– Bob Morris, Sr., former Chief Scientist of the National Computer Security Center, NSA



# Whom or What Am I Trusting? (2)



# Example Secure System Design

- Secure telephone line between FBI and CIA
- Goal: **only people in FBI and CIA buildings can learn what's said in calls**
- Plan:
  - Radiation-proof buildings
  - One entrance/exit per building
  - Armed guards at entrances
  - Guards check ID cards, record all people in/out
  - Pressurized, shielded cable between two buildings
  - No other cables allowed to leave buildings
  - Pass laws to punish people who reveal government secrets
  - Invite NSA to try to steal content of calls
  - Send dummy information, spy on KGB, see if they learn it

# Perfect Security: An Unattainable Goal

- Merely a question of how **motivated** adversary is, and how much **money** he has
- No individual technique perfect
  - Pressurized cable only raises cost for attacker
  - Can't completely shield a building
  - People can be bribed, blackmailed
- Could meet stated goal, but it could be inappropriate
  - What if FBI, CIA allow in uncleared visitors?
  - What if employees go home and talk in sleep?
    - **Solution: forbid employees from leaving the building...**

# Definitions

- **Security:** techniques to control who can access/modify system
- **Principal:** unit of accountability in a system (e.g., user)
- **Access control:** techniques to restrict operations to particular principals
- **Authentication:** verification of identity of principal making request
- **Authorization:** granting of request to principal

# Attacks on Security

- Violation of **secrecy**
  - Attacker **reads data without authorization**
- Violation of **integrity**
  - Attacker **modifies data without authorization**
  - e.g., attacker modifies data on disk
  - e.g., attacker modifies network reply to “read file” request
- **Denial of service**
  - Attacker makes system **unavailable to legitimate users**
  - e.g., overload the system, or cause a deadlock
  - e.g., trigger security mechanism (wrong ATM PIN 3 times)

# Building Secure Systems: General Approach

- Figure out what you want to protect, what it's worth
- Figure out which attacks you want to defend against
- State goals and desired properties clearly
  - Not “impossible to break”
  - Better: “attack X on resource Y should cost \$Z”
- Structure system with two types of components:
  - **Trusted**: must operate as expected, or breach
  - **Untrusted**: subverted operation doesn't lead to breach
- Minimize size of trusted components
  - Maybe we should have built secure **room**, not **building**...
- Analyze resulting system, monitor success

# Security Is a Negative Goal

- Ensure nothing happens without authorization
  - How do you reason about what a system will **not** do?
- First step: specify who authorized to do what
  - In other words, specify a **policy**

# Policy

- Policy: goal security must achieve
  - Human intent—originates from outside system
- Often talked about in terms of subjects and objects
  - Subject: principal
  - Object: abstraction to which access requested (e.g., file, memory page, serial port)
  - Each object supports different kinds of access (e.g., read or write file, change permissions, ...)
- Access control: should operation be allowed?
  - What principal making request? (Authentication)
  - Is operation permitted to principal? (Authorization)



# Access Control: Examples

- Machine in locked room, not on network
  - Policy: only users with keys can access computer
- Bank ATM card
  - Policy: only allowed to withdraw money present in your account
  - Authentication: must have card and know PIN
  - Authorization: database tracks account balances
- Private UNIX file (only owner can read)
  - Authentication: password to run software as user
  - Authorization: kernel checks file's permission bits
- Military classified data
  - If process reads "top-secret" data, cannot write "secret" data

# **Next: User Authentication**

- How to use passwords to authenticate users: at the console, and remotely, over a network
- Attacks against password-based authentication schemes
- Designing robust password-based authentication schemes

# Authentication of Local Users

- Goal: only file's owner can access file
- UNIX authentication policy:
  - Each file has an owner principal: an integer **user ID**
  - Each file has associated **owner permissions** (read, write, execute, &c.)
  - Each **process** runs with integer user ID; only can access file as owner if matches file's owner user ID
  - OS assigns user ID to user's **shell process** at login time, authenticated by **username and password**
  - Shell process creates new child processes with **same user ID**
- How does UNIX know the correspondence among <username, user ID, password>, for all users?

# Straw Man: Plaintext Password Database

- Keep password database in a file, e.g.:  
    bkarp:3715:secretpw  
    mjh:4212:multicast
- Passwords stored in file in **plaintext**
- Make file readable only by privileged **superuser** (root)
- `/bin/login` program prompts for usernames and passwords on console; runs as root, so can read password database
- **How well does this scheme meet original goal?**

# Cryptographic Primitive: Cryptographic Hash Function

- Don't want someone who sees the password database to **learn users' passwords**
- Cryptographic hash function,  $y=H(x)$  such that:
  - $H()$  is **preimage-resistant**: given  $y$ , and with knowledge of  $H()$ , computationally infeasible to recover  $x$
  - $H()$  is **second-preimage-resistant**: given  $y$ , computationally infeasible to find  $x' \neq x$  s.t.  $H(x)=H(x')=y$
- Widely used cryptographic hash functions:
  - MD-5: output is 128 bits, **broken**
  - SHA-1: output is 160 bits; **on verge of being broken**
  - SHA-256: output is 256 bits, **best current practice**

## Better Plan: Hashed Password Database

- Keep password database in a file:

bkar p:3715:Xc8zOP0ZHJkp

mjh:4212:p6FsAtQl4cwi

- Instead of password plaintext  $x$ , store  $H(x)$
- Make file readable by all (!)
- One-wayness of  $H()$  means no one can recover  $x$  from  $H(x)$ , right?
  - **WRONG! Users choose memorable passwords...**

# Insight: Counting Possible Passwords

- If users pick random n-character passwords using c possible characters, how many guesses expected to guess one password?

$$c^n/2$$

e.g., 8 characters, each ~90 possibilities,  **$2.15 \times 10^{15}$**

- **Do users pick random passwords?**
  - Of course not; very hard to remember
  - Common choice: word in native language
- How many words in common use in modern English?
  - **50,000-70,000** (or far fewer, if you read Metro)

# Dictionary Attack on Hashed Password Databases

- Suppose hacker obtains copy of password file (until recently, world-readable on UNIX)
- Compute  $H(x)$  for 50K common words
- String compare resulting hashed words against passwords in file
- **Learn all users' passwords that are common English words after only 50K computations of  $H(x)$ !**
- **Same hashed dictionary works on all password files in world!**